



EFFECTIVE PROGRAMMING TECHNIQUES - less watts

Ivana Hutařová Vařeková, Marcela Mašláňová,
Power Management team, Brno

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1. Design level
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Watt-Hunting competition

- **programmer competition**
- program the maximum of tasks and optimize them in 45 minutes
- use whatever language you want to get best results
- nice awards
- 13:00, 14:00, 15:00, 16:00 in lab B007



1. Design level

- used algorithm and its complexity
- crucial part of optimization
- different complexity classes
- sorting - $O(n \cdot \ln(n))$, $O(n^2)$



2. Implementation level

1. String functions
2. Threads in different languages
3. Memory allocation functions
4. Fsync
5. Wake-ups of disk



2. Implementation level – string functions I

```
char *strcpy(char *dest, const char *src);
```

```
char *strncpy(char *dest, const char *src, size_t n);
```

```
void *memcpy(void *dest, const void *src, size_t n);
```



2. Implementation level – string functions II

```
char *strcpy(char *dest, const char *src);
```

- per byte
- copy character from `src` to `dest`, finish when it finds the first occurrence of character '`\0`'



2. Implementation level – string functions III

```
char *strncpy(char *dest, const char *src,  
              size_t n);
```

- per byte
- tests whether the character is '\0' and the length is n
- if strncpy reaches character '\0' before it writes n chars to dest, it fills the rest of n characters with '\0'
- if strncpy writes n characters before it reaches character '\0' then it stops too – the result string is not end by '\0'



2. Implementation level – string functions IV

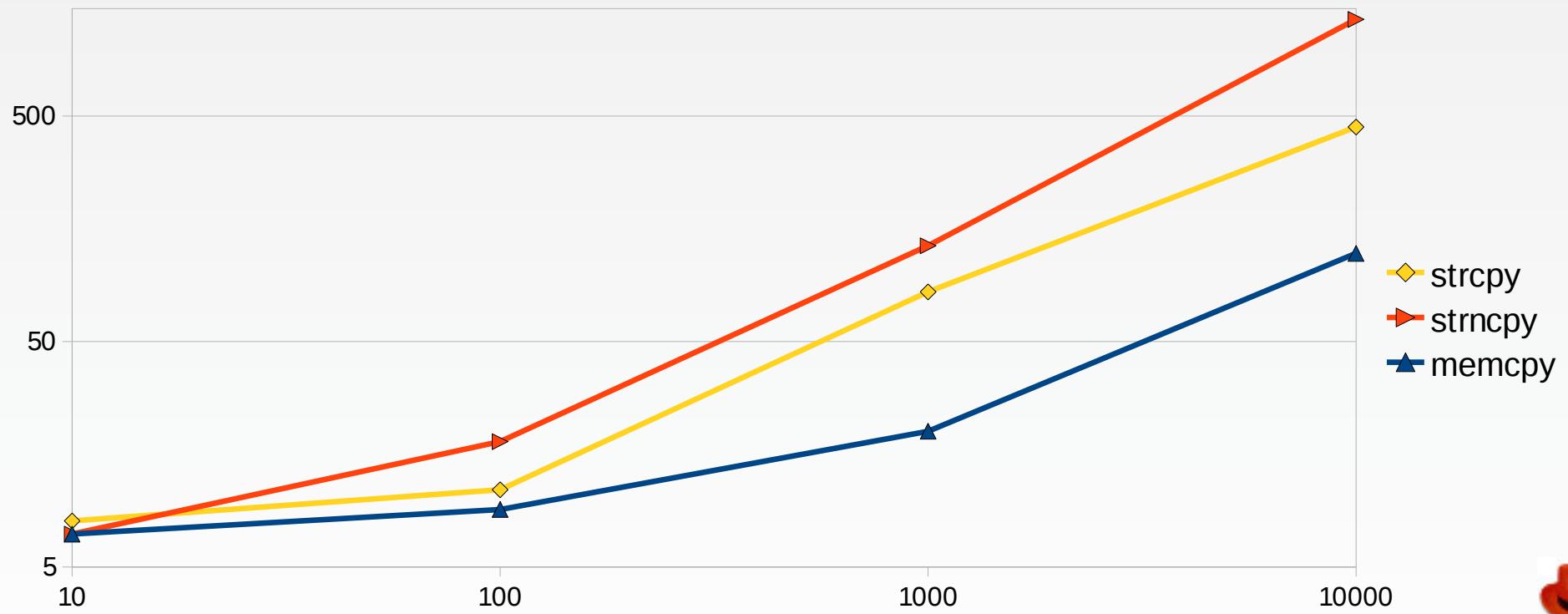
```
char *memcpy(void *dest, const void *src,  
            size_t n);
```

- per word
- copy the memory until it reach the limit of size n



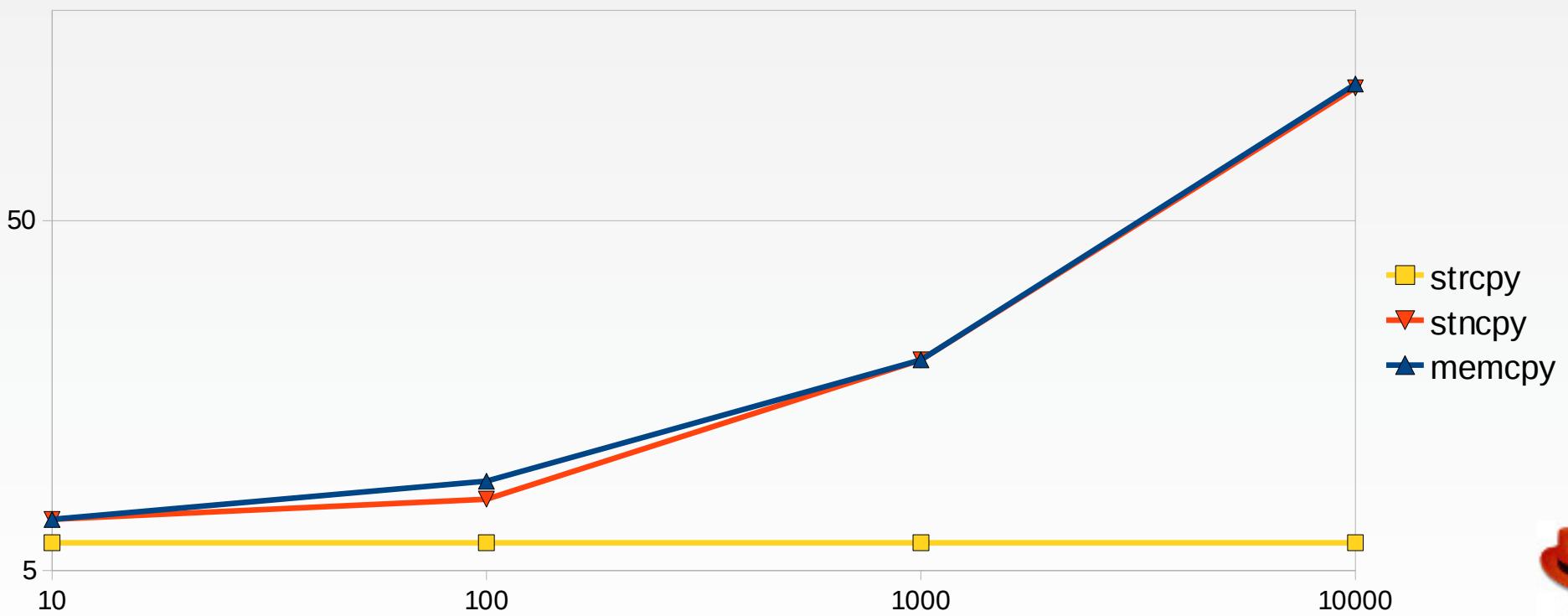
2. Implementation level – string functions V

- `strcpy/strncpy/memcpy` which copies 10/100/1000 character word to buffer of the same size, n is the length of copying word



2. Implementation level – string functions VI

- `strcpy/strncpy/memcpy` which copies 2-character word to 10/100/1000 buffer, n is set to the size of buffer



2. Implementation level – string functions VII

```
char *strcat(char *dest, const char *src);
```

- go through the whole dest character find the first occurrence of '\0'
- then copy the src there
- if we want to copy several strings and use a lot of strcats - the complexity is $O(n^2)$
- if we save the end position and `memcpy` the string there - the complexity is $O(n)$



2. Implementation level – string functions VIII

```
char *stpcpy(char *dest, const char *src);
```

- `stpcpy` – new in POSIX.2008
- copy a string returning a pointer to its end

```
char *to = buf;
```

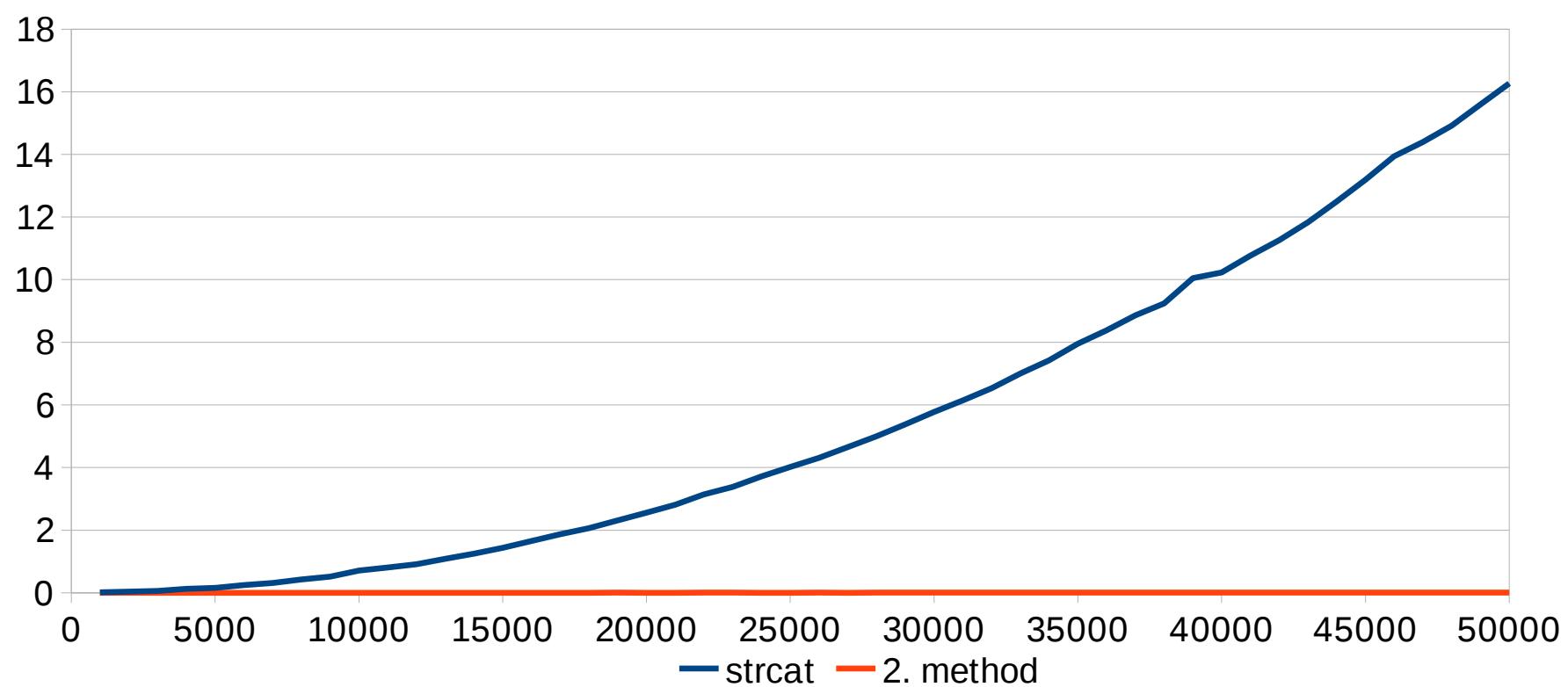
```
to = stpcpy(to,"foo");
```

```
to = stpcpy(to,"bar");
```

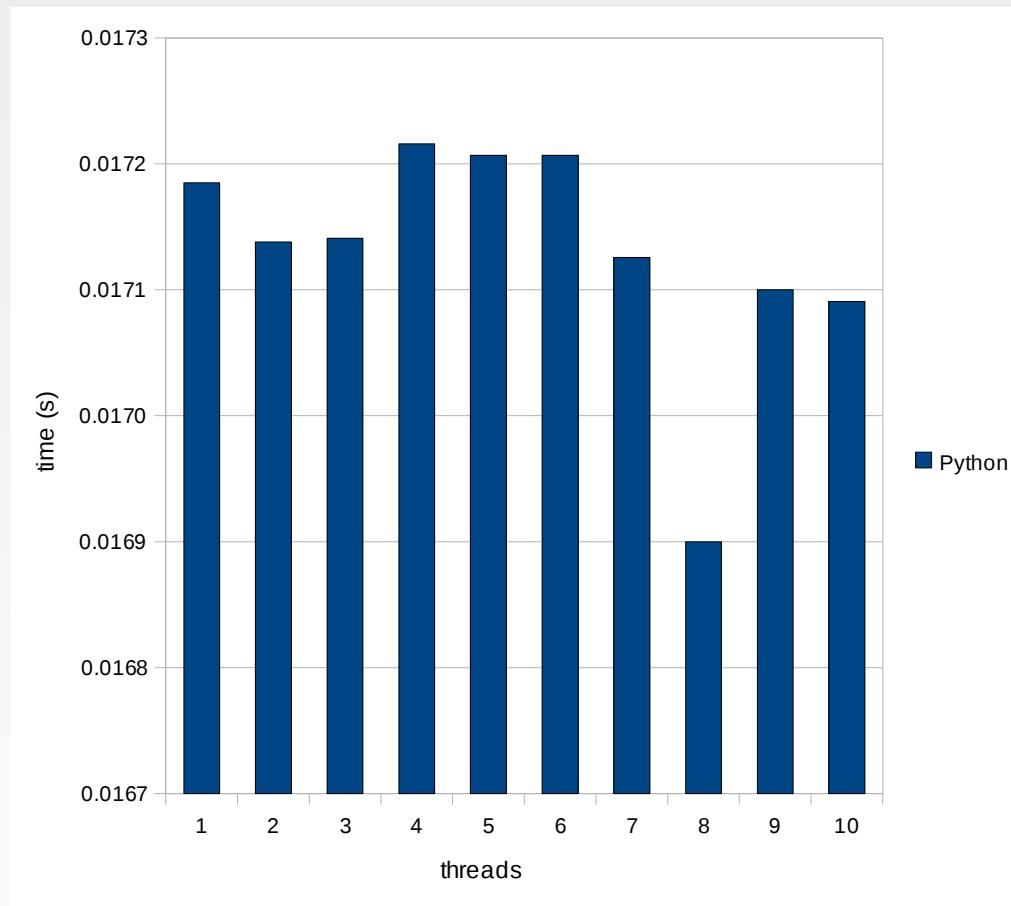


2. Implementation level – string functions IX

cycles	1000	2000
2. method	0.01	0.01
strcat	0.1	0.03



2. Implementation level - threads



Python

- Global Lock Interpreter (GIL)
- operations aren't safe without lock
- profitable only on I/O operations
- optimize with project unladen swallow



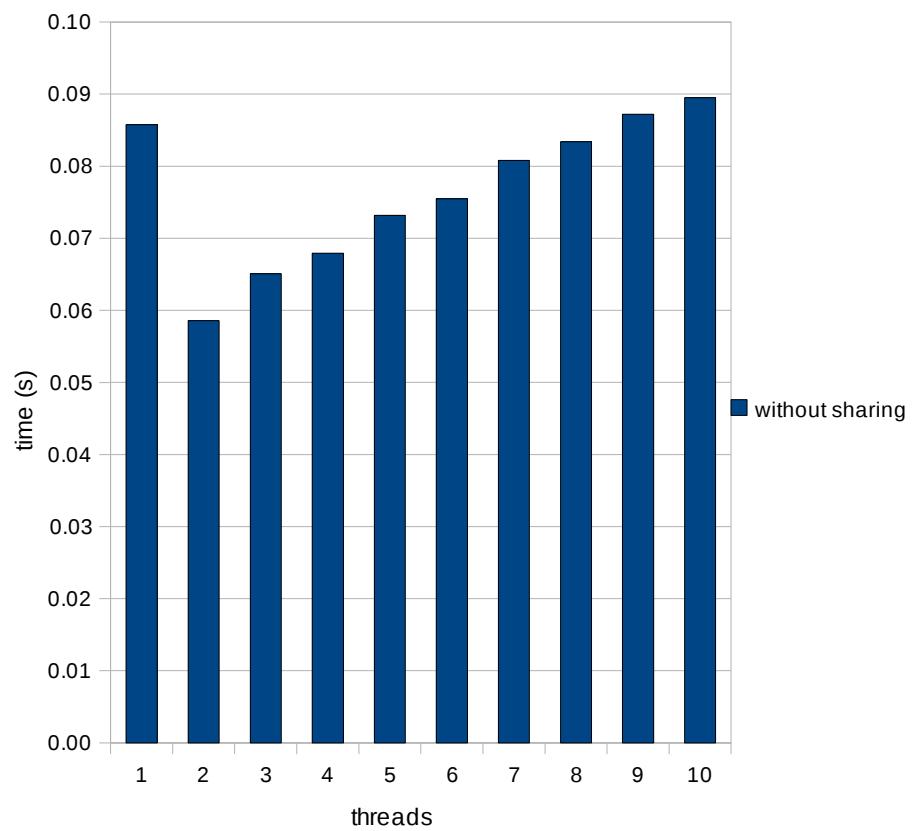
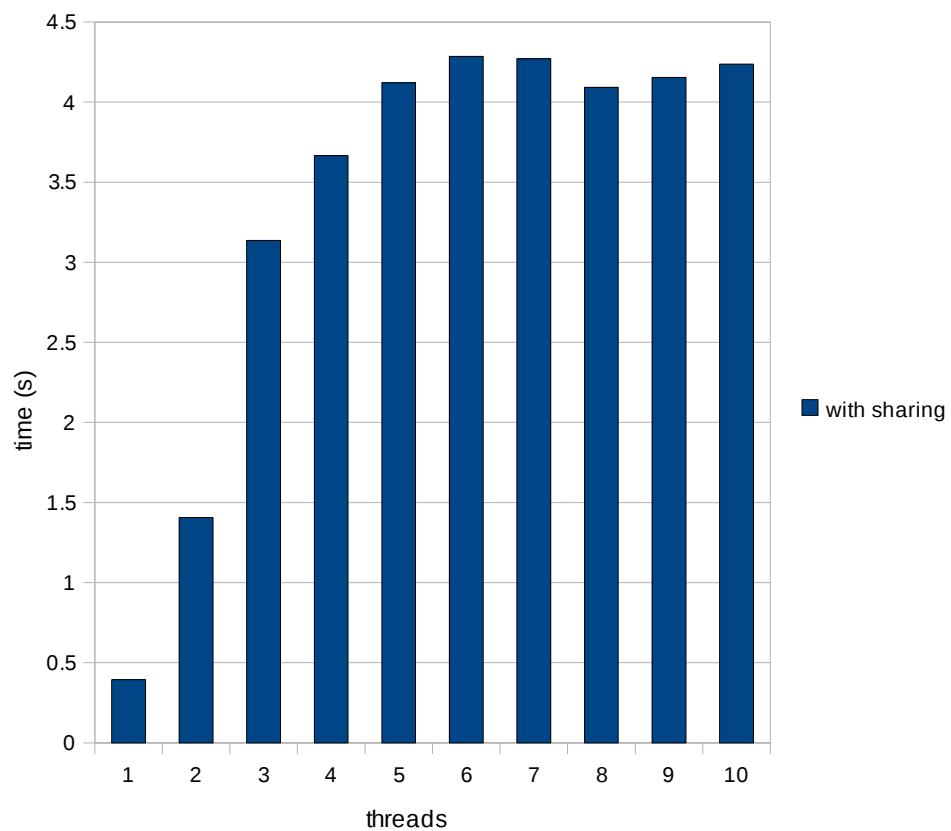
2. Implementation level - threads

- Perl
 - threads, ithreads
 - (don't) share data by default
 - user defines level of sharing
 - Problems:
 - without sharing – COW for every thread
 - with sharing – COW + create tied variables
 - fork could be better for unix only application



2. Implementation level - threads

Multiply matrix 40x40



2. Implementation level - threads

- C
 - threads share the same global memory
 - each thread has his own stack
 - independent threads can access the same variables in memory
 - kernel doesn't create new copy of process memory space and file descriptors
 - standard fork is slower compared to threads
 - threads can take advantage of more CPU

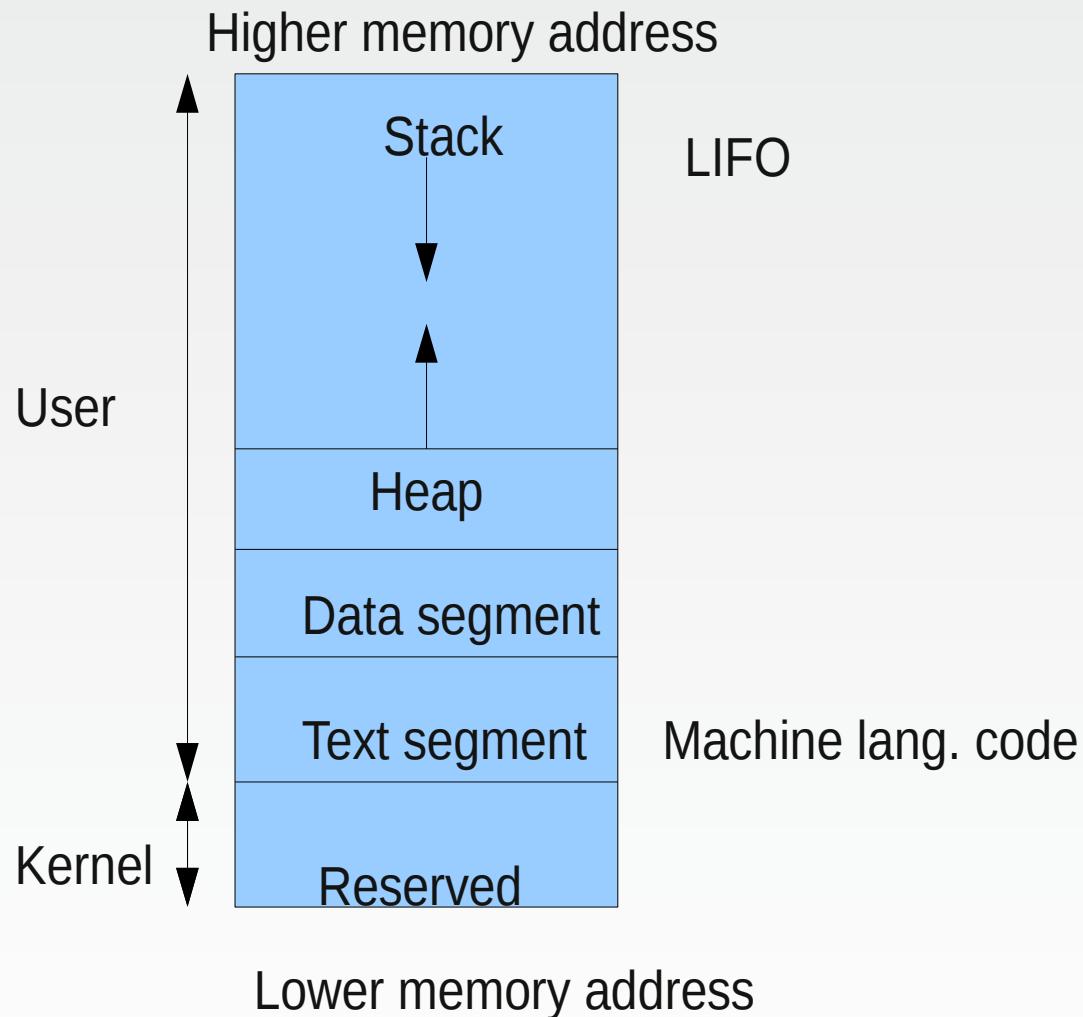


2. Implementation level - memory

- Virtual memory divided into segments:
 - text – code, read-only data
 - data (=heap) – RW segment of binaries/libraries
 - stack
 - segment grows with variables but don't shrink with them
 - pointer at structure is better than structure (time)
 - allocation of larger locale variables dynamically (free'd after end of function/block)



2. Implementation level - memory



2. Implementation level - memory

```
6482 pts/6      Ss+      0:00 bash -rcfile .bashrc
```

```
cat /proc/6482/maps
```

00400000-004d1000 r-xp 00000000 fd:00 39898	/bin/bash
006d0000-006da000 rw-p 000d0000 fd:00 39898	/bin/bash
...	
008d9000-008e1000 rw-p 000d9000 fd:00 39898	/bin/bash
00e9e000-00f01000 rw-p 00000000 00:00 0	[heap]
328ba00000-328ba1d000 r-xp 00000000 fd:00 13883	/lib64/libtinfo.so.5.7
328ba1d000-328bc1d000 ---p 0001d000 fd:00 13883	/lib64/libtinfo.so.5.7
328bc1d000-328bc21000 rw-p 0001d000 fd:00 13883	/lib64/libtinfo.so.5.7
7fa10d08b000-7fa10d08c000 rw-p 00000000 00:00 0	
7fa10d08c000-7fa10d08d000 r--p 0001f000 fd:00 32810	/lib64/ld-2.10.90.so
7fa10d08d000-7fa10d08e000 rw-p 00020000 fd:00 32810	/lib64/ld-2.10.90.so
7fa10d08e000-7fa10d08f000 rw-p 00000000 00:00 0	
7fff4d262000-7fff4d277000 rw-p 00000000 00:00 0	[stack]
7fff4d3ff000-7fff4d400000 r-xp 00000000 00:00 0	[vds0]
ffffffffffff600000-ffffffffffff601000 r-xp 00000000 00:00 0	[vsyscall]



2. Implementation level - memory

- Data segment – allocation functions
 - calloc – allocates and sets the chunk to zero
 - malloc – allocates the chunk
 - realloc – changes size in both directions but increase could be time consuming (not enough memory after allocated chunks)
 - mmap – great for large files, immediately free'd by unmap
 - obstack – pool of memory containing stack objects



2. Implementation level - malloc_trim

- undocumented function, not in POSIX
- useful in specific cases (parsing with libxml2)
- memory isn't free'd after call free -> inefficient memory
- chunks are reserved for application



2. Implementation level – fsync

Bad example:

```
open and read file ~/.kde/myconfig
fd = open("~/kde/myconfig",
O_WRONLY|O_TRUNC|O_CREAT)
write(fd, bufferOfNewData,
sizeof(bufferOfNewData))
close(fd)
```

Better example:

```
open("~/kde/myconfig", O_WRONLY|O_TRUNC|
O_CREAT); read(myconfig);
fd = open("~/kde/myconfig.suffix", O_WRONLY|
O_TRUNC|O_CREAT)
write(fd, bufferOfNewData,
sizeof(bufferOfNewData))
fsync; /* optional - to be 100% sure */
some_operation()
close(fd)
rename("~/kde/myconfig", "~/kde/myconfig~") /*
optional */
rename("~/kde/myconfig.suffix",
 "~/kde/mvconfia")
```

- Call fsync too much
vs. never
- Examples: frozen
Firefox, empty files



2. Implementation level - wake-ups

- wake-ups/spin disk
- bug hunt for power hungry application found some of them
- rewrite checks
 - better design
 - inotify

```
Package gdm
#define
SC_SECURITY_TOKEN_MONITOR_POLL_INTER
VAL 100 /* ms */
```

```
Package hplip rhbz#204725
/* in seconds */
try:
-
    loop(timeout=0.5)
+
    loop(timeout=5.0)
```



2. Implementation level – wake-ups

- inotify watch predefined changes on file/directory since kernel 2.6.13
- since kernel 2.6.31 new API fanotify
- problem: number of watches

/proc/sys/fs/inotify/max_user_watches

```
int fd;
fd = inotify_init();
int wd;
wd = inotify_add_watch(fd,
"../myConfigFile", IN_MODIFY);
...
fd_set rdfs;
struct timeval tv;
int retval;
FD_ZERO(&rfds);
FD_SET(0, &rfds);

tv.tv_sec = 5; /* 5 sec */
value = select(1, &rdfs, NULL, NULL, &tv);
if (value == -1)
    perror("select()");
else {
    do_some_stuff();
}
...
```



3. Compilation level I

gcc optimization flags

- **-O0** - no optimization of result binary, optimized the time of compilation, implicit
- **-O1 (-O)** - optimization which does not want a lot of time and improve the speed of the result binary
- **-O2** - used in fedora flags, all optimization which improve the speed of result binary except of space-speed tradeoff



3. Compilation level II

- **-O3** - all optimizations which improve the speed and can't be problematic
- **-Os** - all optimizations which reduce the size of binary
- **-march=name** - name of the target architecture



4. Profilers (bottleneck) I

- “*We should forget about small efficiencies, say about 97% of the time: premature optimization is the root of all evil. Yet we should not pass up our opportunities in that critical 3%.*”
- Donald Knuth
- “*Bottlenecks occur in surprising places, so don't try to second guess and put in a speed hack until you have proven that's where the bottleneck is.*”
- Rob Pike



4. Profilers (bottleneck) II

- **gprof**
 - needs to have profiling binary and the libraries it need recompiled with -pg flag
- **oprofile**
 - uses kernel module and daemon which have to be start
 - quite complicated managing
- **valgrind --tool=callgrind**



4. Profilers (bottleneck) III

```
valgrind --tool=callgrind [copt] program [poft]
```

- create the data file callgrind.out.\$pid

```
callgrind_annotate callgrind.out.$pid
```

- create the profiling for processes which finished

```
callgrind_control callgrind.out.$pid
```

- create the profiling for processes which are running



4. Profilers (bottleneck) IV

- useful options:

--threshold=<0-100>

set the percentage of counts we are interested in

--tree=none|caller|calling|both
print for each function their callers, the called
function, both

--inclusive=yes|no

add subroutine cost to function calls



4. Profilers (bottleneck) V

1,274,566 PROGRAM TOTALS

Ir file:function

```
1,162,636 * gzip-1.3.12/gzip.c:main [gzip-1.3.12/gzip]
    48 > /usr/src/debug/glibc/signal/sigempty.c:sigemptyset (1x) [/lib/libc-2.9.so]
   180 > /usr/src/debug/glibc/signal/sigaddset.c:sigaddset (6x) [/lib/libc-2.9.so]
   186 > /usr/src/debug/glibc/signal/sigismem.c:sigismember (6x) [/lib/libc-2.9.so]
   902 > gzip-1.3.12/util.c:gzip_base_name (1x) [gzip-1.3.12/gzip]
1,150,531 > gzip-1.3.12/gzip.c:treat_file (1x) [gzip-1.3.12/gzip]

1,150,531 * gzip-1.3.12/gzip.c:treat_file [gzip-1.3.12/gzip]
    77 > gzip-1.3.12/lib/open-safer.c:open_safer (1x) [gzip-1.3.12/gzip]
   835 > gzip-1.3.12/util.c:xunlink (1x) [gzip-1.3.12/gzip]
  2,024 > gzip-1.3.12/gzip.c:open_and_stat (1x) [gzip-1.3.12/gzip]
1,136,108 > gzip-1.3.12/zip.c:zip (1x) [gzip-1.3.12/gzip]
```



4. Profilers (bottleneck) VI

```
1,136,108 * gzip-1.3.12/zip.c:zip [gzip-1.3.12/gzip]
   15 > gzip-1.3.12/util.c:updcrc (1x) [gzip-1.3.12/gzip]
  864 > gzip-1.3.12/util.c:flush_outbuf (1x) [gzip-1.3.12/gzip]
   63 > gzip-1.3.12/util.c:gzip_base_name (1x) [gzip-1.3.12/gzip]
1,024,770 > gzip-1.3.12/deflate.c:deflate (1x) [gzip-1.3.12/gzip]

1,024,770 * gzip-1.3.12/deflate.c:deflate [gzip-1.3.12/gzip]
269,438 > gzip-1.3.12/trees.c:flush_block (1x) [gzip-1.3.12/gzip]
378,069 > gzip-1.3.12/deflate.c:longest_match (2082x) [gzip-1.3.12/gzip]
 98,223 > gzip-1.3.12/trees.c:ct_tally (2443x) [gzip-1.3.12/gzip]
   77 > gzip-1.3.12/deflate.c:fill_window (1x) [gzip-1.3.12/gzip]
```



4. Profilers (bottleneck) VII

1,274,566 PROGRAM TOTALS

Ir file:function

1,162,857	/usr/src/debug/glibc-/csu/libc-start.c:(below main) [/lib/libc-2.9.so]
1,162,636	gzip-1.3.12/gzip.c:main [gzip-1.3.12/gzip]
1,150,531	gzip-1.3.12/gzip.c:treat_file [gzip-1.3.12/gzip]
1,136,108	gzip-1.3.12/zip.c:zip [gzip-1.3.12/gzip]
1,024,770	gzip-1.3.12/deflate.c:deflate [gzip-1.3.12/gzip]
378,069	gzip-1.3.12/deflate.c:longest_match [gzip-1.3.12/gzip]
269,438	gzip-1.3.12/trees.c:flush_block [gzip-1.3.12/gzip]
188,152	gzip-1.3.12/trees.c:compress_block [gzip-1.3.12/gzip]
122,698	gzip-1.3.12/bits.c:send_bits [gzip-1.3.12/gzip]
110,359	/usr/src/debug/glibc-20081113T2206/elf/rtld.c:_dl_start [/lib/ld-2.9.so]
109,068	/usr/src/debug/glibc-20081113T2206/elf/rtld.c:_dl_main [/lib/ld-2.9.so]
98,223	gzip-1.3.12/trees.c:ct_tally [gzip-1.3.12/gzip]



Materials, links

- Ulrich Dreppers slides

<http://people.redhat.com/drepper>

- gcc documentation

<http://gcc.gnu.org/onlinedoc/gcc-4.4.1/gcc>

- profilers documentation

<http://sourceware.org/binutils/docs/gprof>

<http://oprofile.sourceforge.net/doc/index.html>

- valgrind documentation

<http://valgrind.org/docs/manual/cl-manual.html>



Materials, links

- Python threads

<http://docs.python.org/c-api/init.html?highlight=gil>

<http://code.google.com/p/unladen-swallow/>

- Perl threads

http://www.perlmonks.org/?node_id=288022

- fsync problems

<http://thunk.org/tytso/blog/2009/03/15/dont-fear-the-fsync/>

- inotify

<http://www.linuxjournal.com/node/8478/print>

